



ANSWER BOOK

The Simplex algorithm

Decision Mathematics 1 · Year FM

7 questions · 31 marks

NAVY EDITION

SECTION A · Warm-up

- A1** Stop when no negative entries remain in the objective row. A negative entry [2]
means increasing that variable would still improve the objective, so while one remains
the solution can be bettered. Once none remain, no variable can be increased without
reducing the objective, and the current vertex is optimal. B1 for the stopping condition,
B1 for the reason.
- A2** Add slack variables: $x + y + s_1 = 10$ and $2x + 3y + s_2 = 24$, with the objective as P [4]
 $-4x - 5y = 0$. The tableau rows are **P**: 1, -4, -5, 0, 0 | 0; **s₁**: 0, 1, 1, 0 | 10; **s₂**: 0, 2, 3, 0, 1 | 24. B1
for the slack variables, B1 for the objective row, B1 B1 for the two constraint rows.

SECTION B · Standard

- B1** The most negative entry in the objective row is -5, so the pivot column is **y**. The [6]
ratio test gives $10 \div 1 = 10$ and $24 \div 3 = 8$, so the pivot row is the **s₂** row and the pivot
element is 3.
Divide that row by 3 to give the new **y** row **0**, $2/3$, 1, 0, $1/3$ | 8. Subtract it from the **s₁** row
for **0**, $1/3$, 0, 1, $-1/3$ | 2, and add five times it to the objective row for **1**, $-2/3$, 0, 0, $5/3$ | 40.
So $P = 40$ with $y = 8$, $x = 0$ and $s_1 = 2$. B1 for the pivot column, M1 for the ratio test, A1
for the pivot row, M1 for the row operations, A1 for the new **y** row, A1 for the objective row.
The objective row still has $-2/3$ in the **x** column, so a second iteration is needed. Name
the row operations as $R - R(y)$ and $R + 5R(y)$ beside the tableau, since a mark goes on
them.
- B2** No entry in the objective row is negative, so no variable can be increased to [4]
improve **P**: the tableau is **optimal**. The basic variables are $x = 6$ and $y = 4$, the
non-basic slacks are $s_1 = s_2 = 0$, and $P = 44$. B1 for the optimality statement, B1 for $x = 6$,
B1 for $y = 4$, B1 for $P = 44$. Both slacks being zero means both constraints are tight,
matching the graphical answer.
- B3** Ratios: $12 \div 2 = 6$ and $20 \div 4 = 5$. The -1 is skipped entirely: increasing the [3]
entering variable would make that row's value grow rather than shrink, so it places no
limit. The smallest ratio is 5, so the **third** row is the pivot row. M1 for the ratio test, A1
for the third row, B1 for ignoring the negative entry.

SECTION C · Stretch

- C1** The only negative entry left is $-2/3$, so the pivot column is **x**. The ratio test gives [7]
 $2 \div 1/3 = 6$ and $8 \div 2/3 = 12$, so the pivot row is the **s₁** row and the pivot element is $1/3$.
 Multiply that row by 3 to give the new x row **0, 1, 0, 3, -1 | 6**. Subtract two thirds of it from
 the y row for **0, 0, 1, -2, 1 | 4**, and add two thirds of it to the objective row for **1, 0, 0, 2, 1 | 44**.
 No negative entries remain, so the tableau is optimal: $x = 6$, $y = 4$, $s_1 = s_2 = 0$ and $P = 44$.
 B1 for the pivot column, M1 for the ratio test, A1 for the pivot row, M1 for the row
 operations, A1 for the new x row, A1 for the other two rows, A1 for the solution.
 Negative entries away from the objective row are perfectly normal and are ignored by
 the ratio test. Only the pivot column entries matter there, and only the positive ones.
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- C2** With only $\leq$ constraints and non-negative right-hand sides, the origin is [5]
 feasible and the slack variables form a valid starting solution. A $\geq$ constraint with a
 positive right-hand side makes the origin infeasible, and subtracting a surplus variable
 does not replace what the slacks provided: setting the real variables to zero would
 need a negative surplus, so there is no obvious starting basis to read off. An artificial
 variable is added to each such row to supply one, and must then be driven out. The
two-stage method first minimises the sum of the artificial variables, driving them to
 zero and so reaching a feasible corner, then runs ordinary Simplex from there. The
big-M method instead subtracts a large multiple M of each artificial variable from the
 objective, making them so expensive that the algorithm eliminates them on its own. B1
 B1 for why the ordinary start fails, B1 for the two-stage method, B1 for the big-M method,
 B1 for how the artificial variables are driven out.
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